DD2552 - Seminars on Theoretical Computer Science, Programming Languages and Formal Methods, Seminar 4

Karl Palmskog (palmskog@kth.se)

2021-09-09

Last seminar:

- more PCTL operators
- deciding PCTL formulas on DTMCs
- intuitions behind CTMCs and CSL

Today:

- CTMCs and CSL in detail
- examples

 $\mathcal{M} = (S, s_i, Q, L)$ where

- S is a (finite) set of states
- $s_i \in S$ is the initial state
- $Q: S \times S \mapsto R^{\geq 0}$ is a matrix where
 - $Q(s_j, s_k)$ is the rate of transition from s_j to s_k
 - $Q(s_j, s_j)$ is constrained to be $-(\sum_{i \neq k} Q(s_k, s_j))$

• $L: S \mapsto 2^{AP}$

- let $\mathcal{M} = (S, s_i, Q, L)$
- suppose Q(s,s') = r and $t \in R^{>0}$
- probability of ${\cal M}$ moving from s to s' within time t is $1 e^{-rt}$
- slogan: rate is expentially distributed with parameter Q(s, s')
- define $E(s) = \sum_{s' \in S} Q(s, s')$, "total rate"

•
$$S = s_0, s_1, s_2, s_3$$

• $s_i = s_0$
• $L(s_0) = \emptyset, L(s_1) = \{\}, L(s_2) = \{\}, L(s_3) = \{\}$
• $Q(s_0, s_1) = \$$,

$$egin{array}{ll} \phi ::= & op \mid a \mid
eg \phi \mid \phi \wedge \phi \mid P_{\geq heta}(\psi) \ \psi ::= \phi \; U^{\leq t} \; \phi \ t \in R^{\geq 0}, \; \; heta \in [0,1] \end{array}$$

possible additions:

- unbounded until (straightforward)
- next operator (tricky)
- steady state operator (some work)

Paths for CTMCs

- paths π must track the time that was spent in each state
- approach 1: define CTMC path as function from positive reals to states
- approach 2: define CTMC as infinite state sequence with real-valued transition labels

$$\pi = s_0 \rightarrow^{t_0} s_1 \rightarrow^{t_1} \ldots$$

with $t_i > 0$ and $Q(s_i, s_{i+1}) > 0$.

Define:

- $\pi(t)$: state at time t)
- $\pi[i]$: the *i*th state

By "translation" from PCTL:

Tricky ones:

•
$$\mathcal{M}, \pi \models \phi \ U^{\leq t} \ \phi'$$

• $\mathcal{M}, s \models P_{\geq \theta}(\psi)$